

MAT 108

Winter 2021
RIGHT BOARD



Math	English	Prop?
$\sin(\sqrt{1000})$	the brown dog	No no = or verb,
$\sin(\sqrt{1000}) < \frac{1}{2}$	My dog is brown	Yes
$\sin(x) < \frac{1}{2}$	The dog is brown.	No undef. variable
$\sin(x) < 2$	The dog is brown or not brown.	Yes always true.
$P = NP$	Jan 5 2022 will be rainy in Davis.	Yes hard to evaluate
$x = \sqrt{1000}$ and $\sin(x) < \frac{1}{2}$		Yes.
similar possible but tricky.	This sentence is false.	No (paradox.)

In break out rooms:

Find examples of propositions.
and not propositions.

(Maybe a sentence for
which it is not clear.)

Building new propositions from
old ones.

If P and Q are propositions
write $\sim P$ for not P

$P \wedge Q$ for P and Q (A),

$P \vee Q$ for P or Q .

Summarize via a truth table:

P	Q	$\sim P$	$P \wedge Q$	$P \vee Q$	$\sim P \vee Q$	$\sim [P \wedge \sim Q]$
T	T	F	T	T	T	T
T	F	F	F	T	F	F
F	T	T	F	T	T	T
F	F	T	F	F	T	T

call $\sim P \vee Q$
and $\sim [P \wedge \sim Q]$

P

equivalent. Q

~~Eq. equiv.~~ Eng. sentences:
My dog is either not brown or outside.
has propositional form:
 $\sim P \vee Q$

My dog is not both brown and inside.

$$\sim [P \wedge \sim Q]$$

So as above these English sentences
are equivalent.

$$\frac{\sim P \vee Q}{\begin{array}{c} T \\ F \\ T \\ T \end{array}}$$

Names:

$\sim$	not	, negation , denial
$\wedge$	and	, conjunction
$\vee$	or	, disjunction

$\Rightarrow$ implies, conditional, if - then -
 $\Leftrightarrow$ iff, if and only if, biconditional

$P \Rightarrow Q$ means if $\underbrace{P}_{\text{antecedent}}$ then $\underbrace{Q}_{\text{consequent}}$.

converse to $P \Rightarrow Q$ is $Q \Rightarrow P$

contrapositive to $P \Rightarrow Q$ is $(\neg Q) \Rightarrow (\neg P)$.

Break out rooms!

(Even) Find 2 Eng. sentences
involving conditions which
have to ^{mean} some meaning but
diff. prop. forms

(Odd) Find two diff equiv. prop.
forms in v. condit. and then

Find Eng. sentences
modeling the form.

Ex: $P \Rightarrow Q$

[If I'm inside I am at Zoom U.]

$\sim Q \Rightarrow \sim P$ (contrapos.)

If I'm not at Zoom I'm outside.

Ex: I'm a student at UC Davis P

I'm a math major Q.

$$\left. \begin{array}{l} P \Rightarrow Q \\ Q \Rightarrow P \end{array} \right] \text{ clearly not equ.}$$

$P \Leftrightarrow Q$ $\xleftrightarrow{\text{equiv.}}$ $(P \Rightarrow Q) \wedge (Q \Rightarrow P)$

T
T
T
T
T

T
T
T
T
T

Rules for getting equivalent propositional forms:

Thm 1.1.1 $\sim \sim P \quad \text{eq} \quad P$

$$P \wedge Q \quad \text{eq} \quad Q \wedge P$$

$$P \wedge (Q \wedge R) \quad \text{eq} \quad (P \wedge Q) \wedge R$$

$$P \wedge (Q \vee R) \quad \text{eq} \quad (P \wedge Q) \vee (P \wedge R)$$

De Morgan $\sim (P \vee Q) \quad \text{eq} \quad (\sim P) \wedge (\sim Q)$

$$\sim (P \wedge Q) \quad \text{eq} \quad (\sim P) \vee (\sim Q)$$

Ex: Find a prop. form ~~for~~ equiv to

$\sim(P \Rightarrow Q)$ without conditionals
($\Rightarrow, \Leftrightarrow$).

Ans: Truth Table:

PQ		$\sim(P \Rightarrow Q)$	$P \wedge \sim Q$
T	T	F	F
T	F	T	T
F	T	F	F
F	F	F	F

Ans:

$$P \wedge \sim Q$$

Another Ans:

$P \Rightarrow Q$ is eq. to

$$\sim P \vee Q$$

(from before).

so $\sim(P \Rightarrow Q)$ eq. $\sim[\sim P \vee Q]$

eg $P \wedge \sim Q$

210108

Here a universe (or discourse) is
the set over which all ~~quant~~
quantifiers are taken.

(eg: $\forall x$ means $\forall x \in \text{universe}$)

Problem: With a paradox like
this sentence is false

Consider the denial of this ~~statement~~
sentence:

Somebody I know dislikes chocolate
~~No body~~ and likes coffee.

$$(\exists x \in \{\text{ppl I know}\}) (\sim P(x) \wedge Q(x))$$

Not everyone I know likes chocolate
or dislikes coffee.

$$\sim (\forall x \in \{\text{ppl I know}\}) (P(x) \vee \sim Q(x))$$

$\exists$ is there exists:

$$(\exists x \in \mathbb{Z}) (x^2 = 1)$$

Means there is some integer x
with $x^2 = 1$

or
in the universe $\mathbb{Z}$.

$$(\exists x) (x^2 = 1).$$

In the universes below which are true?

U	$P(x)$ all integers = $\{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$ " $\mathbb{Z}$	$\mathbb{R}$	$\mathbb{R}_{\geq 0}$	$\{-\sqrt{5}, \sqrt{5}\}$
	$\mathbb{Z}$	$\mathbb{R}$	$\mathbb{R}_{\geq 0}$	$\{-\sqrt{5}, \sqrt{5}\}$
$(\forall x) P(x)$	F	F	F	T
$(\exists x) P(x)$	F	T	T	T
$(\exists! x) P(x)$	F	F	T	F

Notation: $\exists!$ is there is exactly one.
(there exists a unique).

Proof of Thm 1.3.1-a:

Let U be any universe.

The sentence $\neg (\forall x) P(x)$ is true in U
 $(\exists x) P(x)$ is ~~false~~ in U

$\xrightarrow{\text{iff}}$
if and only if

iff the truth set for P in U
is not all of U

iff the truth set for $\sim P$ in U
is not empty.

iff $(\exists x) \sim P(x)$

Thus $\sim (\forall x) P(x)$ is eqv. to $(\exists x) \sim P(x)$
in any universe -

210111

Thm: If a is an odd integer
then a^2+1 is an even integer.

Proof: Assume that a is an odd integer.
Hence there is an integer n with
 $a = 2n+1$.

$$\begin{aligned}\text{Hence } a^2+1 &= (2n+1)^2+1 = 4n^2+4n+1+1 \\ &= 2[2n^2+2n+1] \\ &\text{and } 2n^2+2n+1 \text{ is an integer.}\end{aligned}$$

There some a^2+1 is an even integer.

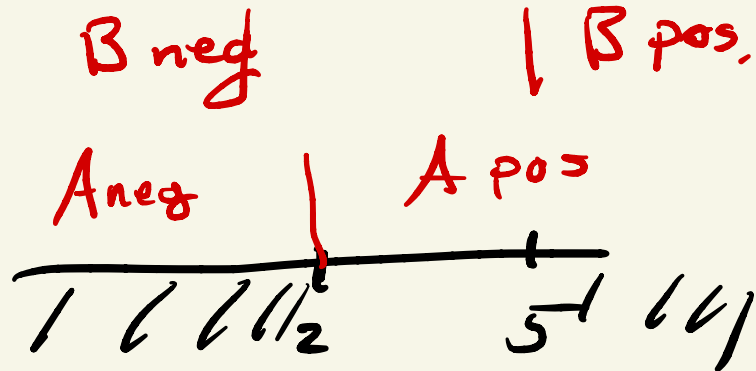
g.ed. latin sor: that which was
to be proved

Why is it true?

work backward and forward:

back

$$\left\{ \begin{array}{l} x^2 - 7x > -10 \\ \text{or} \\ x^2 - 7x + 10 > 0 \\ \text{or} \end{array} \right.$$



Ex 2.17

$(x-2)(x-5) > 0$

A or B

A & B are both pos or A & B are both neg

or $x < 2$ or $x > 5$

$$x^2 \leq 1 \quad \text{or} \quad -1 \leq x \leq 1$$

hence $x < 2$ ✓

Proof Plan: Show $x^2 \leq 1$ implies $x < 2$

then hence $x-2 < 0$ & $x-5 < 0$

then $(x-2)(x-5) > 0$

done $x^2 - 7x > -10$

Try a similar approach to show.

Thm: If x is a real number with

$x < 1$ then $x^2 + 6 > 5x$.

Proof:

Assume $x < 1$.

Hence $x < 2$.

Hence $(x-3)(x-2) > 0$.

Hence $x^2 + 6 > 5x$.

qed.

210113 Another direct proof;

Def: If a and b are integers
then a divides b if
 $(\exists n \in \mathbb{Z})$ with $n \cdot a = b$

Claim: If a, b and c are integers
and a divides both b and c
then a divides $b - c$.

Proof: Assume a, b and c are integers
with a dividing b and c .

Hence there ~~is~~ are integers n and m
with $n \cdot a = b$ and $m \cdot a = c$.

Hence $n - m$ is an integer.

Hence, $b - c = n \cdot a - m \cdot a = (n - m) \cdot a$.

Therefore a divides $b - c$. qed.

Understood!

eg $5 = m$ $5^2 = 25$ odd so 5 odd ✓
 $4 = m$ $4^2 = 16$ even so 4 even ✓

Logic: $(\forall m \in \mathbb{Z}) \quad P \Rightarrow Q$

P is m^2 is odd

Q is m is odd. *use this.*

or equivalently: $\sim Q \Rightarrow \sim P$

A: Assume: $\sim Q$ or m is even

or $\exists t$ with $2 \cdot t = m$

Conclusion: $\neg P$ or m^2 is even

or $\exists s$ with $2 \cdot s = m^2$

Ideas: $m^2 = (2t)^2 = 4t^2 = 2(\underbrace{2t^2})$ ✓
an int.

Claim: Assume a is an integer.

4 does not divide $a^2 - 2$.

210115

Recall §1.7 read often.

claim ① getting started;

Understand: e.g.: $6^2 - 2 = 34 = 2 \cdot 17$ not a mult of 4

Logic: P is 4 divides $a^2 - 2$

$$(\forall a \in \mathbb{Z}) \sim P$$

or for contradiction equiv:

$$\sim [(\forall a \in \mathbb{Z}) \sim P] \Rightarrow \underbrace{[Q \wedge \sim Q]}_F \text{ still get to choose } Q.$$

$$\boxed{\text{e.g. } [(\exists a \in \mathbb{Z}) P] \Rightarrow [Q \wedge \sim Q]}$$

Ass/Concl:

Assume there is some int. a with 4 dividing $a^2 - 2$ (2) ←
or $(\exists t \in \mathbb{Z}) \quad 4t = a^2 - 2$

Concl: $\mathbb{Q}$ and $\sqrt{2}$ still get to change $\mathbb{Q}$,

Idea 1: $4t = a^2 - 2$

(3)

so $4t + 2 = 2(2t + 1) = a^2$

so a^2 is even so by Claim 1 a is even

so there is $s \in \mathbb{Z}$ with $2s = a$

so $4t = (2s)^2 - 2 = 4s^2 - 2$

so $2 = 4(s^2 - t)$

⑤ $1 \in \mathbb{Q}$ so

$$1 = 2(5^2 - 4)$$

an integer

④

so

1 is even

$\mathbb{Q}$

⑥ Know $1 \in \mathbb{Q}$

1 is not even

$\sim \mathbb{Q}$

Quote! §1.7 Go over your proof and

- check for things not needed
- Make sure the order is best
- Make sure it is complete.

Def: Write $a \in \mathbb{Q}$ or a is rational
if $\exists p \in \mathbb{Z}, \exists q \in \mathbb{N}$ with $a = \frac{p}{q}$.

Claim ③: If a is rational $\forall a \in \mathbb{Q}$

and $a = \frac{p}{q}$ with $p \in \mathbb{Z}$ and $q \in \mathbb{N}$
and the denominator q as small as possible
then either p is odd or q is odd.

U: $a = \frac{2^4}{132} = \frac{12}{66} = \frac{2}{11} \leftarrow \text{odd}$ ✓

$a = \frac{3}{5} \leftarrow \text{odd}$ ✓
 $a = \frac{3}{10}$ ✓

$$L: (\forall a \in \mathbb{Q}) \underbrace{P}_{\text{complicated}} \Rightarrow \underbrace{Q}_{\text{simple}} \vee \underbrace{R}_{\text{simple} \rightarrow}$$

or contrapositive:

$$(\forall a \in \mathbb{Q}) \sim(Q \vee R) \Rightarrow \sim P$$

or

$$(\forall a \in \mathbb{Q}) (\sim Q \wedge \sim R) \Rightarrow \sim P$$

Thm (Pythagoreans ~ 230 BC)
 $\sqrt{2} \notin \mathbb{Q}$

or $\sqrt{2}$ is not rational

$$r \neq \frac{7}{5}$$

Claim ④ There is an integer n
with $n^2 + 15 < 8n$

Proof: Take $n = 4$.

$$\text{Hence } n^2 + 15 = 31$$

$$\text{and } 8n = 32 > 31.$$

qed.

210120)

so

$$2q^2 = 2(2m+1)^2 = 8m^2 + 8m + 2$$

$$p^2 = (2k)^2 = 4k^2$$

so

$$4m^2 + 4m + 1 = 2k^2$$

so

$$1 = 2k^2 - 4m^2 - 4m$$

$$= 2[k^2 - 2m^2 - 2m]$$

1 is even

Q

8 ✓

know

1 is not even

~ Q

9

From §1.7 (read often)

Always use complete sentences in a proof.

Multiple quantifiers:

- ① $\exists!$ both $\exists$ and $\forall$ something else.
- ② $\forall - \exists -$ or $\exists - \forall -$.

Example: 3 possibly false statements:

Claim: $\sim [(\exists! x \in \mathbb{R}) \quad x^2 - 4x + 5 = 0]$

Proof: It suffices to show:

(sketch) $\sim [(\exists x \in \mathbb{R}) \quad x^2 - 4x + 5 = 0]$

Computing: $x^2 - 4x + 5 = (x^2 - 4x + 4) + 1$

$$= (x-2)^2 + 1$$

For contradiction: assume $x \in \mathbb{R}$ and

$$x^2 - 4x + 5 = 0 \quad \text{so} \quad (x-2)^2 + 1 = 0$$

$$\text{so} \quad (x-2)^2 = -1 < 0$$

Know $(x-2)^2 \geq 0$

Then we have $(x-2)^2 \geq 0$ and $(x-2)^2 < 0$

a contradiction so $\neg (\exists x \in \mathbb{R}) x^2 - 4x + 5 = 0$

QED

210122

Ex:	a	b	$\gcd(a, b)$	$\text{lcm}(a, b)$
	3	18	3	18
	10	35	5	70
	91	91	91	91
	256	81	1	$256 \cdot 81$

Def: (pg 82 ex 16)

If a, b and m are nonzero integers
then $\text{lcm}(a, b) = m$ iff

(i) a and b both divide m
(say m is a common multiple of
 a and b)

(ii) every common multiple of a and b
is at least m .

$(\forall a, b, m \in \mathbb{Z} \neq 0) \text{lcm}(a, b) = m \iff$

$$\begin{aligned}
 & [(\exists s, t \in \mathbb{Z}) (a \cdot s = m) \wedge (b \cdot t = m)] \wedge \\
 & [(\forall e \in \mathbb{Z}) [(\exists u, v \in \mathbb{Z}) (a \cdot u = e) \wedge (b \cdot v = e)] \\
 & \Rightarrow (e \equiv m)]
 \end{aligned}$$

Shortening strategy:

Def: $\underbrace{a \mid e} \Leftrightarrow (\exists s \in \mathbb{Z}) a \cdot s = e$
 a divides e

Def: Call a and b relatively prime

or coprime if $(a, b) = 1$.

Approach to (a):

u: $a = 20$ $b = 30$ / $\gcd(20, 30) \neq 10$

$$b =$$

$$a = 11$$

$$b = 13$$

$$\gcd(11, 13) = 1 \wedge 1 \nmid 13$$

$$L: (\forall a)(\exists b) P \wedge Q$$

$$P \text{ is } P_i \wedge P_{ii}$$

$P_i \text{ is } 1 \text{ divides both}$

always true

P_{ii} and comm div- is at most 1

A/c: $\forall$.. assume $a \in \mathbb{N}_{\geq 0}$
find (our choice) b
conclude: $P \wedge Q$.

I: eg $a = 20$
choose $b = 23$ ✓

if: $a = 65$
choose: $b = 83$ ✓

Try:
 a
 $b = a + 1$

If d is a common div of
 a and $a+1$

$\exists s, t \in \mathbb{Z}$ with

$$d \cdot s = a \quad \text{and} \quad d \cdot t = a+1$$

$$\text{so} \quad d \cdot t - d \cdot s = a+1 - a = 1$$

$$\parallel \\ d(t-s)$$

210125

Ex: $\left[\begin{array}{l} 2 \text{ is a lin comb of 10 and 14.} \\ \text{since: } 2 = 10 \cdot 3 + 14 \cdot (-2) \end{array} \right.$

$\left[\begin{array}{l} 1 \text{ is not a lin. comb. of 10 and 14} \\ \text{since: For contradiction assume} \\ 1 = 10 \cdot x + 14 \cdot y = 2 \cdot [5 \cdot x + 7 \cdot y] \\ \text{so 1 is odd and 2 is even. a contradiction.} \end{array} \right.$

Ex of 1.8.1:

The smallest positive linear combination
of 10 and 14 is 2.

Also $\gcd(10, 14) = 2$.

Ex: If $a=10$, $b=14$ then

$$14 = 10 \cdot 1 + 4 \quad \text{and} \quad 0 \leq 4 < 10$$

Euclid's Alg. Example: From $a=10$, $b=14$

$$14 = 10 \cdot 1 + 4$$

$$10 = 4 \cdot 2 + \boxed{2} = r_k$$

$$4 = 2 \cdot 2$$

No

┌

$$q_1 = 1$$

$$q_2 = 2$$

$$q_3 = 2$$

$$k = 2$$

$$r_k = 2$$

$$256 = 81 \cdot 3 + 13$$

$$81 = 13 \cdot 6 + 3$$

$$13 = 3 \cdot 4 + \boxed{1} \quad r_2 = \gcd(256, 81),$$

$$3 = 1 \cdot 3$$

U: $p = 3$

$p = 6$
not true

$$a = 4$$

$$b = 9$$

$$a = 4$$

$$b = 9$$

$$ab = 36$$

$p \nmid ab$ but not a or b

p is prime must be used.

$$L: (\forall a, b, p \in \mathbb{Z}) (P \wedge Q) \Rightarrow (S \vee T)$$

P is p is prime

Q is $p \mid ab$

S is $p \mid a$

T is $p \mid b$

divides.

$\sim S \Rightarrow T$

equivalently:

$$\neg (\forall a, b, p \in \mathbb{Z}) (P \wedge Q \wedge \sim S) \Rightarrow \neg T$$

Ass: $P, Q, \sim 5$

Cond: T

I: $\left[\begin{array}{l} p \text{ is prime} \\ p \mid ab \\ p \nmid a \end{array} \right] \begin{array}{l} \text{only divis are 1 and } p \\ \text{so } ab = pn \text{ for some} \\ \text{int. } n. \\ \text{so } \gcd(p, a) = 1 \end{array}$

① $\underbrace{p \nmid a}_{\text{does not divide}}$

By
Thm: 1.9.1 $\left[\begin{array}{l} \text{② so } 1 = p \cdot x + a \cdot y \\ \text{for some ints } x \text{ and } y. \\ \text{gcd}(p, a) = \text{lin. comb. of } p \text{ and } a \end{array} \right]$

$\exists x, y \in \mathbb{Z}$

Playing with these:

$$ab = pn$$

$$[1 = p \cdot x + a \cdot y$$

(4)

want

$$b \stackrel{?}{=} p \cdot t$$

$$\rightarrow b = p \cdot x \cdot b + a \cdot y \cdot b$$

$$= p \cdot xb + p \cdot n \cdot y$$

$$= p[xb + ny]$$

so

$p \text{ div } b.$

$$\begin{aligned} & (\forall p, a \in \mathbb{Z} \neq 0) \\ & (\exists x, y) \quad \text{gcd}(a, b) \\ & \quad = a \cdot x + b \cdot y \end{aligned}$$

(T).

(5)



210127 Chapter 2: Set Theory -

Notes ① Will consider set theory using first order logic.

② Also any first order logic involves sets

- Ⓐ the universe U is a set.

- Ⓑ If $P(x)$ is a proposition with variable in a universe U then the truth set of P is $\{x \in U \mid P(x) \text{ (is true)}\}$

Consider relationships between operations on sentences and on sets.

Def: If A and B are sets then A is a subset of B if

$$(\forall x)[(x \in A) \Rightarrow (x \in B)]$$

In this case write $A \subseteq B$.



Lemma: If A is a set then

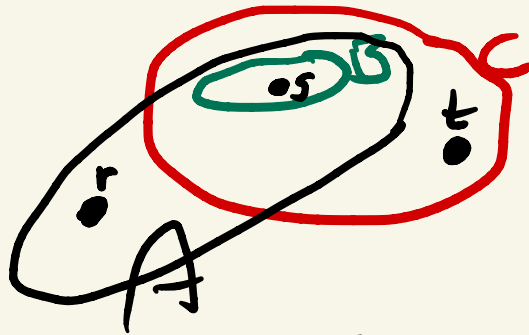
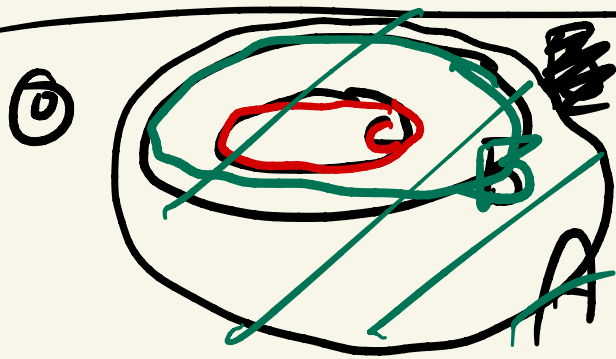
$$\{\} \subseteq A.$$

Proof: $(\forall x) (x \neq \{\})$

So $(x \in \{\})$ is false

and implies $(x \in A)$.

Then Reflex: $(\forall x) [(x \in \{\}) \Rightarrow (x \in A)]$
qed.



True: Example: $A = \{r, s\}$

$$B = \{s\}$$

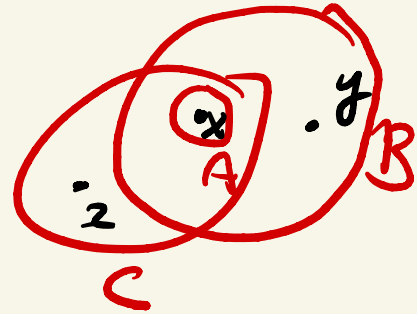
$$C = \{s, t\}$$

check ✓,

$$\textcircled{1} \quad A = \{x\}$$

$$B = \{x, y\}$$

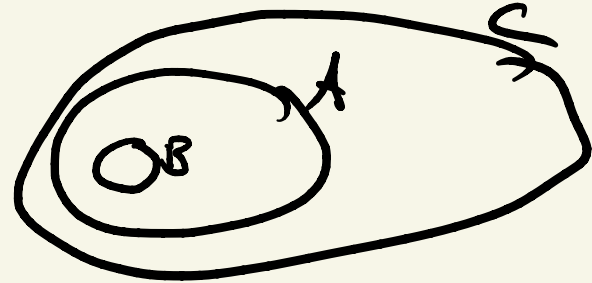
$$C = \{x, z\}$$



$$\textcircled{2} \quad A = \{x, y\}$$

$$B = \{x, y\}$$

$$C = \{x, y\}$$



③

④ False : Proof idea:

Know:

$$B \subseteq A \quad \text{so}$$

$$(\forall x) (\text{~~if~~ } x \in B) \Rightarrow (x \in A)$$

$$A \subseteq C \quad \text{so}$$

$$(\forall x) (x \in A) \Rightarrow (x \in C)$$

$$\text{so } (\forall x) (x \in B) \Rightarrow (x \in C).$$

so

$$B \subseteq C$$

210129

Claim: $\{x \in \mathbb{R} \mid |x+3| = 5 - |x|\} = \{-4, 1\}$.

PS plan: ① check: 1 in $\rightarrow$ try $|1+3| \stackrel{?}{=} 5 - |1| = 4 \checkmark$.

② Need to show: $A \subseteq B$ and $B \subseteq A$.

③ Assume $x \in A$ and show $x \in B$

• A... $x \in B$... $x \in A$

④: $A = LHS = \{x \in \mathbb{R} \mid \dots\}$ $B = RHS = \{-4, 1\}$

To show $B \subseteq A$ just plug both in.

To show $A \subseteq B$

$|x|$ suggests cases:

$x < -3$: $-(x+3) = 5+x$ or $2x = -8$ or $x = -4$ ✓

$-3 < x < 0$: $x+3 = 5+x$ or $3=5$ no soln

$0 < x$: $x+3 = 5-x$ or $2x=2$ ✓

$$2x = -8 \quad x = -4$$

Proof: Check: $\{-4, 1\} \subseteq \{x \in \mathbb{R} \mid |x+3| = 5 - |x|\}$

If $x = 1$ then $|x+3| = 5 - |x| = 4$ ✓.

If $x = -4$ then $|x+3| = 5 - |-4| = 1$ ✓

Check: $\{-4, 1\} \supseteq \{x \in \mathbb{R} \mid |x+3| = 5 - |x|\}$ qed.

$\overset{B}{\parallel} \overset{A}{\parallel}$

Assume $x \in A$. Either $0 < x$ or $-3 \leq x \leq 0$
or $x < -3$.

If $0 < x$ then $x+3 = 5-x$ so $x = 1 \in B$

If $-3 \leq x \leq 0$ then $x+3 = 5+x$ so $3 = 5$
a contradiction. so $x \in B$.

If $x < -3$ then $-x-3 = 5-x$ so $x = -4 \in B$.

Therefore $x \in B$. qed

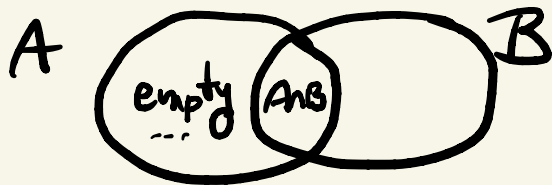
Problem: Which is true?

① If A & B are sets then
 $A \subseteq B$ iff $A \cap B = A$.

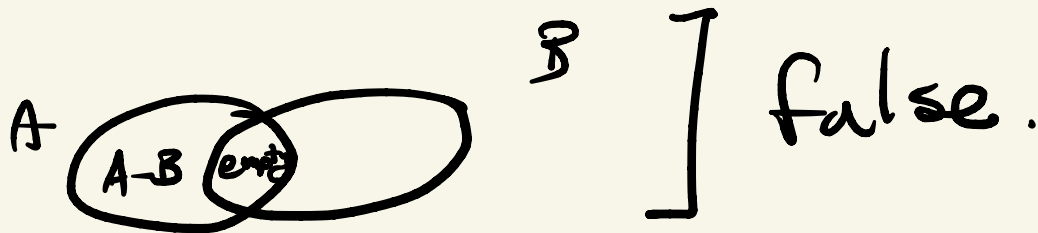
② If A & B are sets then
 $A \subseteq B$ iff $A - B = A$.

Ans: Approach:

①



②



For ② Find a counter example:

Set $A = \{1\}$ $B = \{2\}$.

compute: $A - B = \{1\} = A$ } RHS true

but $A \not\subseteq B$ } LHS false.

210261

$$A = \{2, 3\}$$

$$B = \{3\}$$

$$C = \{4\}$$

$$D = \{5\}$$

$$A \times B = \{(1, 3), (2, 3)\}$$

$$C \times D = \{(4, 5)\}$$

$$(A \cup C) \times (B \cup D) = \{(1, 3), (2, 3), (4, 3), (1, 5), (2, 5), (4, 5)\}$$

so
but $A \times B = C \times D = \{\}$
 $(A \cup C) \times (B \cup D) = \{(4, 3)\}$

Asside: $\{\} \times \{3\} = \{\}$

Note: For any set B
 $\{\} \times B = \{\}$ $(0 \cdot b = 0) \checkmark$

Ex: $\{\{\}\} \times \{3\} = \{(\{\}, 3)\}$

Note: If A has a elements
and B has b elements
then $A \times B$ has $a \cdot b$ elements.

Note: $\{\} \subseteq B$ for any set B .

Proof that $(\forall A, B, C, D \text{ sets})$

$$(A \times B) \cup (C \times D) \subseteq (A \cup C) \times (B \cup D)$$

Assume: (x, y) is in $A \times B$

so x is in A and y is in B .

Hence since $A \subseteq A \cup C$ and $B \subseteq B \cup D$
also have $x \in A \cup C$ and $y \in B \cup D$.

Therefore (x, y) is in $(A \cup C) \times (B \cup D)$.

If (x, y) is in $C \times D$ then similarly,
 (x, y) is in $(A \cup C) \times (B \cup D)$.

QED

Def: If $\mathcal{Q}$ is a family of sets.

write $\bigcup_{A \in \mathcal{Q}} A$ for the union

so $x \in \bigcup_{A \in \mathcal{Q}} A$ iff $(\exists A \in \mathcal{Q}) x \in A$

If $\mathcal{Q} = \{A_\alpha \mid \alpha \in \Delta\}$ is an indexed family
of sets then $x \in \bigcup_{\alpha \in \Delta} A_\alpha$ iff

$$(\exists \alpha \in \Delta) x \in A_\alpha$$

Ex: ①: $\bigcup_{A \in \mathcal{A}} A = \bigcup_{\alpha \in \{a, b, c\}} A_\alpha = \{1, 2, 3, 4, 5\}$

② $\bigcup_{B \in \mathcal{B}} B = \bigcup_{\alpha \in [0, 2)} B_\alpha = [0, 5)$

Def: $x \in \bigcap_{A \in \mathcal{A}} A \iff (\forall A \in \mathcal{A}) \quad x \in A$
 $x \in \bigcap_{\alpha \in \Delta} A_\alpha \iff (\forall \alpha \in \Delta) \quad x \in A_\alpha$

210203

Proof Planning

U:

$$n=1$$

$$1^2 = 1$$

✓

$$n=2$$

$$2^2 = 1 + 3$$

✓

$$n=3$$

$$3^2 = 1 + 3 + 5$$

✓

L: $(\forall n \in \mathbb{N})$

$P(n)$

] correct logic for induction.

I: To show the formula for $n=4$

$$\begin{aligned}
 \underbrace{4^2}_{(n+1)^2} &= \underbrace{(3+1)^2}_{(n+1)^2} = \underbrace{3^2 + 2 \cdot 3 + 1}_{(1+3+5)} + (2 \cdot 3 + 1) \\
 &= 1 + 3 + 5 + \underbrace{7}_{2 \cdot 4 - 1} \quad \checkmark
 \end{aligned}$$

Or more generally:

$$\begin{aligned}
 (n+1)^2 &= (n+1)^2 = \underbrace{n^2 + 2n + 1}_{(1+3+\dots+(2n-1))} + (2n+1) \\
 &= 1 + 3 + \dots + (2(n+1)-1)
 \end{aligned}$$

Got $P(n+1)$ using $P(n)$

Claim: For every natural number n
it is true that $n+3 < 5n^2$.

Proof: check the base case ^{with} $n=1$
which is $1+3 < 5 \cdot 1^2$ which is true.

Assume for induction that $n+3 < 5n^2$

$$\text{hence } 5(n+1)^2 = 5n^2 + 1 + 10n + 4$$

$$> 5n^2 + 1 > n+3+1 = (n+1)+3.$$

Therefore by PMI the claim holds.

q.e.d.

Bad Proof:

Claim: Any group of n horses
are all the same color.

"Proof" (look for error).

Check the base case with $n=1$
which is that any one horse
is one color and is true,

Assume for induction that

any group of n horses
are all the same color.

Hence in any group of $n+1$
horses the first n are all
the same color and the last n
are all the same color and
these two groups both contain the
middle horses so all $n+1$
horses are all the same color.

Then by PMI the claim holds.
QED



Proof Plan:

Consider this an inductive
statement: $(\forall n \in \mathbb{N}) \quad \nexists a \in \mathbb{Z}$

with $\frac{a}{n} = \sqrt{2}$

base case with $n=1$

is $\sqrt{2} \neq \frac{a}{1}$ an integer;

since $0^2 \neq 2$, $1^2 \neq 2$

and if $|a| \geq 2$ then $a^2 > 2$ ✓.

Assume for induction that
for any $n \leq m$ $\nexists \frac{a}{n} = \sqrt{2}$.

Hence: Assume for contradiction
that $\exists b \in \mathbb{Z}$ with $\frac{b}{m+1} = \sqrt{2}$.

so $b^2 = 2(m+1)^2$ so b^2 is even

so $b = 2k$ is even

so $4k^2 = 2(m+1)^2$

so $2k^2 = (m+1)^2$ so $(m+1)^2$ is even

so $m+1 = 2r$ is even

so $\frac{k}{r} = \frac{\frac{2k}{2r}}{m+1} = \frac{b}{m+1} = \sqrt{2}$

but $\frac{m+1}{2} \leq r$ contradicting the
induction hypothesis.

216205

Proof sketch for Claim ①:

U: Note $\gcd(1,1)=1$
 $\gcd(1,2)=1$
 $\gcd(2,3)=1$
 $\gcd(3,5)=1$
 $\vdots$ ✓

L: $(\forall n \in \mathbb{N}) \quad P(n)$

with $P(n)$ being: $\gcd(f_n, f_{n+1}) = 1$

C/A: Use induction:

• Check $n=1$ base case. ②

• Assume $P(n)$ and show $P(n+1)$ (3)
PMI

(Assume $\forall k \leq n$ have $P(k)$)
and show $P(n+1)$
GPMI

I:

$$\begin{aligned} \gcd(13, 8) &= \gcd(8+5, 8) \\ &\stackrel{?}{=} \gcd(5, 8) = 1 \end{aligned}$$

$\underbrace{\hspace{15em}}_{P(5)}$

More generally: $\gcd(f_{n+1}, f_{n+2}) = \gcd(f_{n+1}, f_{n+1} + f_n)$

$$\stackrel{?}{=} \gcd(f_{n+1}, f_n) \stackrel{=}{=} 1$$

ind hyp ✓

Claim (2): If f_n are the Fibonacci numbers then f_{4n} is divisible by 3.

Sketch for (2):

u_i	f_1	f_2	f_3	f_4 3	f_5	f_6	f_7	f_8 21
	1	1	2	3	5	8	13	21

3 div

3 div

L: $(\forall n \in \mathbb{N}) P(n)$

here : $P(n)$ is 3 divides f_{4n}

A/c: Use induction:

base case $n=1$

$(f_4=3)$ ✓

Assume for induction $P(n)$
conclude $P(n+1)$.

I:
$$f_{4(n+1)} = f_{4n+4} = f_{4n+3} + f_{4n+2}$$

$$= f_{4n+2} + f_{4n+1} + f_{4n+2}$$

$$= 2f_{4n+2} + f_{4n+1}$$

$$= 2(f_{4n+1} + f_{4n}) + f_{4n+1}$$

$$= 3f_{4n+1} + 2f_{4n}$$

and by the ind. hyp. $f_{4n} = 3 \cdot 5$

$$\text{so } f_{4(n+1)} = 3(f_{4n+1} + 2 \cdot 5).$$

210208

Ex: § 2.3 : Set families proof;

Prove that if $\mathcal{Q} = \{A_\alpha \mid \alpha \in \Delta\}$
(10a) and $\Gamma \subseteq \Delta$.

$$\text{then } \bigcup_{\alpha \in \Gamma} A_\alpha \subseteq \bigcup_{\alpha \in \Delta} A_\alpha$$

Proof: Assume $a \in \bigcup_{\alpha \in \Gamma} A_\alpha$.

Hence $(\exists \alpha \in \Gamma)$ with $a \in A_\alpha$.

and since $\Gamma \subseteq \Delta$ we have $d \in \Delta$.

There fore $a \in \bigcup_{d \in \Delta} A_d$. qed.

Thm 2.5.2: (Division Algorithm):

If $0 < m \leq n$ ^{are} n natural numbers then
there are unique q and r with

$$n = q \cdot m + r \quad \text{and} \quad 0 \leq r < m$$

Proof of uniqueness:

Assume $0 < m \leq n$ and $n = q \cdot m + r$

$$\text{and } n = \tilde{q} \cdot m + r \text{ and } 0 \leq r < m$$

$$\text{and } 0 \leq \tilde{r} < m,$$

$$\text{Hence: } -m < \tilde{r} - r < m.$$

$$\text{Also } 0 = n - n = (q - \tilde{q}) \cdot m + (r - \tilde{r})$$

$$\text{Hence } \tilde{r} - r = (q - \tilde{q}) \cdot m$$

$$\text{so } \frac{\tilde{r} - r}{m} = q - \tilde{q}.$$

$$\text{Since } -1 < \frac{\tilde{r} - r}{m} < 1 \quad \text{is an integer} \\ \text{we have } \frac{\tilde{r} - r}{m} = 0.$$

$$\text{Hence } \tilde{r} - r = 0 = (q - \tilde{q}) \cdot m.$$

$$\text{Therefore } \tilde{r} = r \text{ and } \tilde{q} = q. \quad \text{QED.}$$

210212 Notation around Relations:

Def: If A and B are sets
then a relation from A to B
is a subset $R \subseteq A \times B$.

write: $a R b$ or a is related to b
if $(a, b) \in R$
 $a \not R b$ or a is not related to b
if $(a, b) \notin R$

The domain of a relation R
from A to B is

$$\text{Dom}(R) = \{a \in A \mid (\exists b \in B) a R b\}$$

and the range is

$$\text{Rng}(R) = \{b \in B \mid (\exists a \in A) a R b\}$$

A relation from A to A is a relation on A .

Ex: The identity relation on A is $I_A = \{(a, a) \in A \times A\}$.

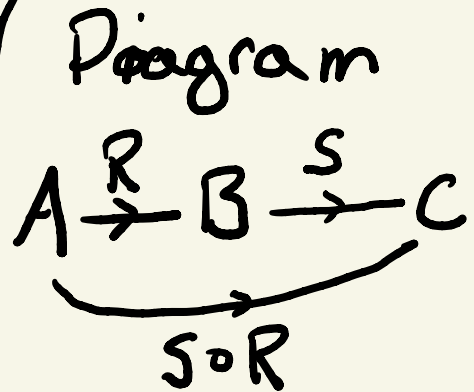
The inverse to a relation R

from A to B is the relation
 R^{-1} from B to A with
 $R^{-1} = \{(b, a) \mid (a, b) \in R\}.$

Note: $I_A^{-1} = I_A$

Def: If R is a relation
from A to B and
 S is a relation
from B to C

Then the composition relation

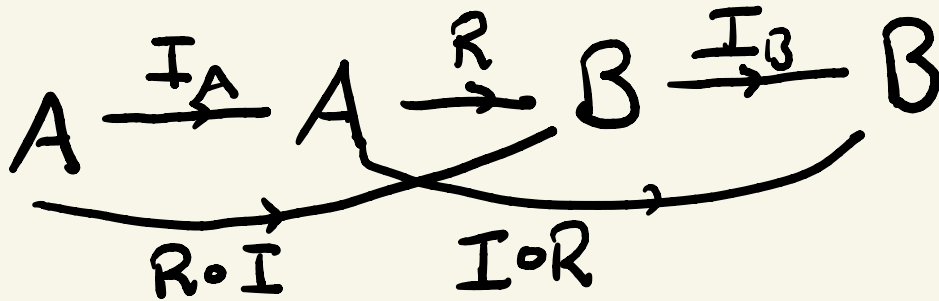


$$S \circ R = \{ (a, c) \in A \times C \mid (\exists b \in B) (a R b) \wedge (b S c) \}$$

from A to C

$$\text{Ex: } \textcircled{a} I_B \circ R = R$$

$$\textcircled{b} R \circ I_A = R$$



Pf: Sketch:

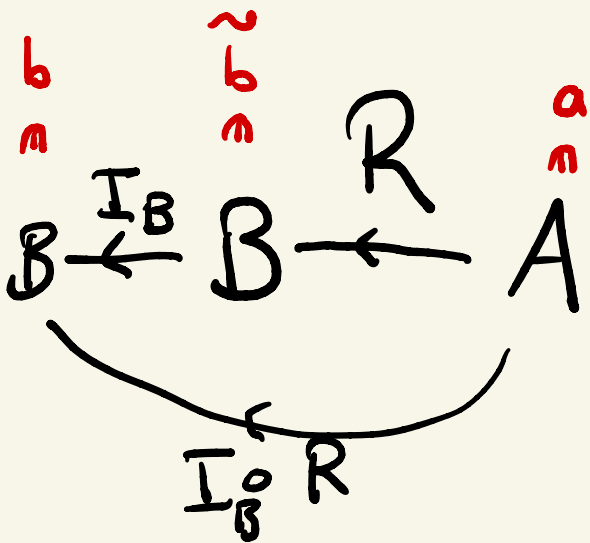
Undertand: R and $I_B \circ R$

Logic: are relations and hence sets
so we must show $I_B R \subseteq R$ and
 $I_B \circ R \supseteq R$

Ass/Concl: • Assume $(a,b) \in I_B \circ R$
and show $(a,b) \in R$.
/ • Assume $(a,b) \in R$
and show $(a,b) \in I_B \circ R$.

Idea: $\hookrightarrow$ If $(a, b) \in I_B \circ R$

then $(a, b) \in I_B \circ R = \{ (a, b) \in A \times B \mid$
 $(\exists \tilde{b} \in B) (a R \tilde{b}) \wedge (\tilde{b} I_B b) \}$.



but $\tilde{b} I_B b$

iff $b = \tilde{b}$

so $(a R \tilde{b}) \wedge (\tilde{b} I_B b)$
 iff $a R b$ and $b = \tilde{b}$

210217

Claim: -4 is not in the range of T

$$\begin{aligned}\text{Proof: } \text{Rng}(T) &= \{y \mid (\exists x \in \mathbb{R}) \ x T y\} \\ &= \{y \mid (\exists x \in \mathbb{R}) \ x^2 = y\}.\end{aligned}$$

Assume for contradiction that $-4 \in \text{Rng } T$.

Hence there is $x \in \mathbb{R}$ with $x^2 = -4 < 0$.

If $x \in \mathbb{R}$ then $x^2 \geq 0$.

Hence $x^2 \geq 0$ and $x^2 < 0$ a contradiction.

Therefore $-4 \notin \text{Rng } T$.

Claim: If R is a relation from A to B .

then $\text{Dom}(R^{-1}) = \text{Rng}(R)$.

Proof: $\text{Rng}(R) = \{y \mid (\exists x \in A) \ x R y\}$
 $= \{y \mid (\exists x \in A) \ y R^{-1} x\}$
 $= \text{Dom}(R^{-1}).$ QED

Ex: $R =$ 

refl: Yes (all loops are true).

Sym: No: $\mathbb{R}^2$ but $\mathbb{Z}$
 trans: Yes. check for every pair
 of composable arrows.

Ex: $U =$

 $= \{ (x, y) \mid x \leq y \}$

refl: Yes: $\{ (x, x) \mid x \in \mathbb{R} \} \subseteq U$
 sym: No: $\circ U$ but $\nexists$
 trans: Yes: if x, y, z with
 $x \leq y$ $y \leq z$ then $x \leq z$.

that is if $x, y, z \in \mathbb{R}$.

and $x \cup y$ and $y \cup z$

then $x \leq y$ and $y \leq z$

so $x \leq z$

so $x \cup z$.

$$Ex: V = \text{graph with a red X} = \{(x, y) \mid x = y \text{ or } x = -y\}.$$

refl:
sym:
~~trans:~~

Y
Y

since for no (x, x) is in V .
since the graph is preserved
by reflection over $x = y$
or!

if $x \vee y$

then either $y=x$ or $y=-x$

if $y=x$ then $y \vee x$.

if $y=-x$ then $y \vee x$

~~done~~

trans: γ

$$V = \{(x, y) \mid x=y \text{ or } x=-y\}$$

$$= \{(x, y) \mid x^2 = y^2\}$$

so if

x, y, z with $x \vee y, y \vee z$

$$\text{so } x^2 = y^2 \text{ and } y^2 = z^2$$

$$\text{so } x^2 = z^2 \text{ so } x \sim z.$$

Hence $\sim$ is an equivalence relation

Sketch:

Logic: Need $\Rightarrow$ and $\Leftarrow$] for iff.

$\Rightarrow$ Assume R is sym.
and show: $\subseteq$
 $\supseteq$

$\Leftarrow$ Assume $R=R^{-1}$
and show symmetric.

210219

If R is a relation on a set A

define:

if $x \in A$ write $\bar{x} = \{y \in A \mid x R y\}$.

Write: $A/R = \{\bar{x} \mid x \in A\}$

a family of subsets of A .

Def: If A is a non empty set
and $\mathcal{F}$ is a family of subsets of A
then $\mathcal{F}$ is a partition of A if

(i) $\emptyset \notin \mathcal{F}$

(ii) If $E, F \in \mathcal{F}$ then $[(E = F) \text{ or } (E \cap F = \emptyset)]$

(iii) $\bigcup_{E \in \mathcal{F}} E = A.$

Thm 3.3.1: If R is an equivalence relation on a nonempty set A then A/R is a partition of A .

²⁶⁰²⁰²
Ex: 7.a:

Find R an equiv. reln on $\mathbb{N}$

so that $\mathbb{N}/R = \{ \{1, 2, \dots, 9\}, \{10, \dots, 99\}, \{100, \dots, 999\}, \dots \}$

Ans: $R = \{ (a, b) \mid \text{if } a \text{ and } b \text{ have the same number of digits} \}$

§3.4: Mod. Arith.

Consider for any $m \in \mathbb{N}$

there is an equiv. reln: on $\mathbb{Z}$

$$R_m = \{(a, b) \mid m \text{ divides } a - b\}.$$

$$\text{Write } \mathbb{Z}_m = \mathbb{Z} / R_m = \{ \underbrace{\overline{0}, \overline{1}, \overline{2}, \overline{3}, \dots, \overline{m-1}}_{\parallel} \}$$

$$\{ \dots, -2m, -m, 0, m, 2m, \dots \}$$

$$\parallel \\ \overline{-1}$$

and call this set the integers mod m .

Def: if $\bar{a}, \bar{b} \in \mathbb{Z}_m$, $k \in \mathbb{N}$ then:

$$\bar{a} + \bar{b} = \overline{a+b} \quad \text{write} \quad a+b \quad (\text{mod } m)$$

$$\bar{a} \cdot \bar{b} = \overline{a \cdot b} \quad a \cdot b \quad (\text{mod } m)$$

$$\bar{a}^k = \overline{a^k} \quad a^k \quad (\text{mod } m)$$

Claim: $+$ is well defined in $\mathbb{Z}_m$.

Equivalently: a is congruent mod m to c

Claim: If $a \equiv c \pmod{m}$ and $b \equiv d \pmod{m}$

then $a+b \equiv c+d \pmod{m}$

Proof: Since m divides $a-c$ and $b-d$

there are integers k and l with

$$mk = a - c \quad \text{and} \quad ml = b - d$$

$$\text{so} \quad mk + ml = (a+b) - (c+d)$$

$$\quad \quad \quad \parallel$$
$$m(k+l)$$

$$\text{so} \quad m \text{ divides } (a+b) - (c+d)$$

$$\text{so} \quad a+b \equiv c+d \pmod{m}$$

qed.

210224

Def: If A and B are sets then
a function f from A to B
is a relation from A to B

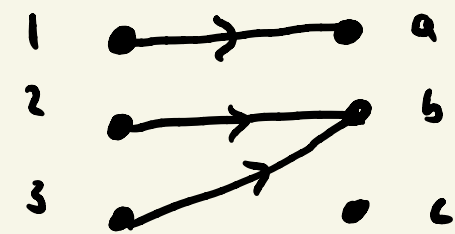
with (i) $\text{domain}(f) = A$

and (ii) if $a \in A$ and $b, c \in B$
with $(a, b) \in f$ and $(a, c) \in f$
then $b = c$.

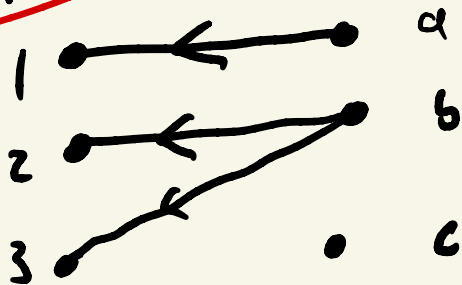
Equiv:

A relation f from A to B is
 a function iff
 $(\forall a \in A) (\exists! b \in B) (a, b) \in f.$

Ex: directed graph:



reln from $\{1, 2, 3\}$
 to $\{a, b, c\}$



reln from $\{a, b, c\}$
 to $\{1, 2, 3\}$

range
 ||
 $\{a, b\}$

codomain
 could be
 $\{a, b\}$
 $\{a, b, c\}$
 $\{a, b, c, \dots\}$

Claim: Example ① is a function:

$$f = \{ (x, y) \in \mathbb{R}^2 \mid y = x^2 - 4 \}.$$

Notation for functions:

If f is a function from A to B
write: $(a, b) \in f$ or equiv $a f b$

or equiv $f(a) = b$

A codomain for a function is any set
containing the range.

Sketch of proof:

U: Note: vertical line test holds ✓

L: (i) For domain: $\forall x \in \mathbb{R}$ have $\exists y$ $f(x) = y$.

(ii) for second (uniqueness)

condition: $\forall f(x) = y \quad \& \quad f(x) = z$
have $y = z$.

I: (i) $\forall x$ choose $y = x^2 - 4$ ✓

(ii) if $x^2 - 4 = y \quad \& \quad x^2 - 4 = z$
then $y = z$ ✓.

Ex: If A is any set
then $I_A = \{(a, a) \mid a \in A\}$ is
a function from A to A
called the identity on A .

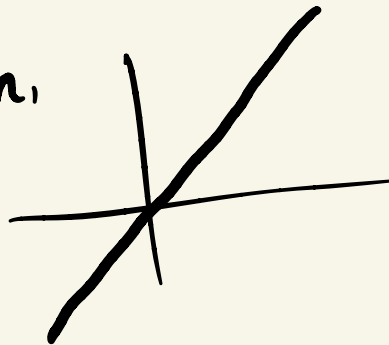
Note: This relation is also an equivalence
relation,

Sketch: U: See above I_A is both,

L: Need to show $f \subseteq I_A$
Ass: f is a fn
and eqv reln. and $f \supseteq I_A$.

I: $\supseteq$ since f is an eq-reln
 f is reflexive
so $\forall a \in A$ here $(a,a) \in f$.

$\subseteq$ since f is a function,
if $f(a) = b$ then $b = c$
and $f(a) = c$



but $f(a) = a$ since $I_A \subseteq f$.

An attempt to define a function using a relation which is not one is called not well defined.

Sketch:

U: The relation is

$$f = \{(\bar{x}, [x]) \mid x \in \mathbb{Z}\}.$$

$$= \{(\bar{0}, [0]), (\bar{1}, [1]), \dots\}.$$

L: Want to show f is not a function.

and $\left\{ \begin{array}{l} \text{(i)} \quad \forall a \in \mathbb{Z}_3 \quad \exists b \in \mathbb{Z}_6 \text{ with } (a,b) \in f. \\ \text{(ii)} \quad \forall a \in \mathbb{Z}_3, b, c \in \mathbb{Z}_6 \text{ with } (a,b), (a,c) \in f. \\ \quad \quad \text{have } b = c \end{array} \right.$

$\sim \textcircled{i}$ or $\sim \textcircled{ii}$.

I: $\sim \textcircled{ii}$ is likely the problem:

$$\{(\bar{0}, [0]) \dots, (\bar{3}, [3])\}$$

$$\begin{array}{c} \parallel \quad \neq \\ \bar{0} \quad [0] \end{array}$$

$\therefore \sim \textcircled{ii}$ is $\exists a \in \mathbb{Z}_3, b, c \in \mathbb{Z}_6$ with $(a, b), (a, c) \in f$
 $\sim (b=c)$
 $(b \neq c),$

Recall: $\bar{x}, \bar{y} \in \mathbb{Z}_3$
then $\bar{x} = \bar{y}$ if 3 divides $x - y$.

$[x], [y] \in \mathbb{Z}_6$
 then $[x] = [y] \iff 6 \text{ divides } x-y$

eg: $\bar{0}, \bar{3} \in \mathbb{Z}_3$ and 3 divides $0-3 = -3$ ✓

Sketch:

w: $g = \{([x], \bar{x}) \mid x \in \mathbb{Z}\} = \{([0], \bar{0}), \dots\}$

Note: $g = f^{-1}$ the inverse relation.

L: (i) $\forall a \in \mathbb{Z}_6 \quad \exists b \in \mathbb{Z}_3$ with
 $(a, b) \in g$

(ii) $\forall a \in \mathbb{Z}_6, \quad b, c \in \mathbb{Z}_3$ with
 $(a, b), (a, c) \in g$ have $b = c$.

I: (i) $\forall a \in \mathbb{Z}_6 \quad \exists n \in \mathbb{Z}$ with $a = [n]$
and $([n], \bar{n}) \in g$ so take $b = \bar{n}$ v.

(ii) If $a \in \mathbb{Z}_6$ and $b, c \in \mathbb{Z}_3$

with $(a, b), (a, c) \in g$

so there is $n \in \mathbb{Z}$ with $a = [n]$

and $k, l \in \mathbb{Z}$ with $b = \bar{k}$
 $c = \bar{l}$

and $([k], \bar{k}) = (a, b)$
 $([l], \bar{l}) = (a, c)$

so $[k] = [l]$

so 6 divides $k - l$ so $k - l = 6t$
 $= 3 \cdot (2t)$

so 3 divides $k - l$ so $\bar{k} = \bar{l}$ so $b = c$.

More Functions:

Characteristic functions;

If S is a subset of A

then $\chi_S = \{(s, 1) \mid s \in S\} \cup \{(a, 0) \mid a \in A - S\}$

is a function from A to $\{0, 1\}$.

In particular $\chi_S(s) = 1$ if $s \in S$

$\chi_S(a) = 0$ if $a \in A$ and $a \notin S$

Inclusion Functions;

If S is a subset of A

the $i_S = \{(s, s) \mid s \in S\}$.

from S to A .

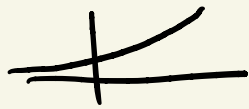
Restrictions of functions:

If S is a subset of A
and f is a function from A to B

then $f|_S = \{(s, f(s)) \mid s \in S\}$

from S to B

Example 5:



Consider: $f(x) = e^x$

from $\mathbb{R}$ to $\mathbb{R}$ or $\mathbb{R}_{>0}$

from $\mathbb{R}$ to $\mathbb{R}$

and $g(x) = x+3$

Find: $f^{-1}(x) = \ln(x)$

from $\mathbb{R}_{>0}$ to $\mathbb{R}$

from $\mathbb{R}$ to $\mathbb{R}$

$g^{-1}(x) = x-3$

from $\mathbb{R}$ to $\mathbb{R}_{>0}$

$(f \circ g)(x) = e^{x+3}$

$(g \circ f)(x) = e^x + 3$

from $\mathbb{R}$ to $\mathbb{R}$

or $\mathbb{R}$ to $\mathbb{R}_{>0}$

4.2-1
Claim: If g & f are functions
as above then $g \circ f$ is a function.

Claim 4.21: If g is a function from B to C and f is a function from A to B then $g \circ f$ is a function from A to C .

Proof: For existence:

If $a \in A$ then $f(a) = b$ is in B so $g(b)$ is in C and (a, c) is in $g \circ f$.

For uniqueness:

If (a, c) and $(a, \tilde{c})$ are in $g \circ f$ then there are b and $\tilde{b}$ with

(a, b) and $(a, \tilde{b})$ in f so $b = \tilde{b}$

and (b, c) and $(b, \tilde{c})$ in g so $c = \tilde{c}$.
qed.

§4.3:

Def: A function f from A to B
is onto or surjective if $\text{range}(f) = B$.

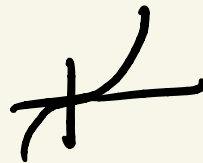
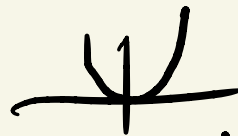
A function g from A to B
is one to one or injective if
every x and y in A with $f(x) = f(y)$
has $x = y$.

Examples: Which of the following
are onto and/or one to one
as functions from $\mathbb{R}$ to $\mathbb{R}$?

① $x^2 = f(x)$

② $e^x = g(x)$

③ $x^3 = h(x)$

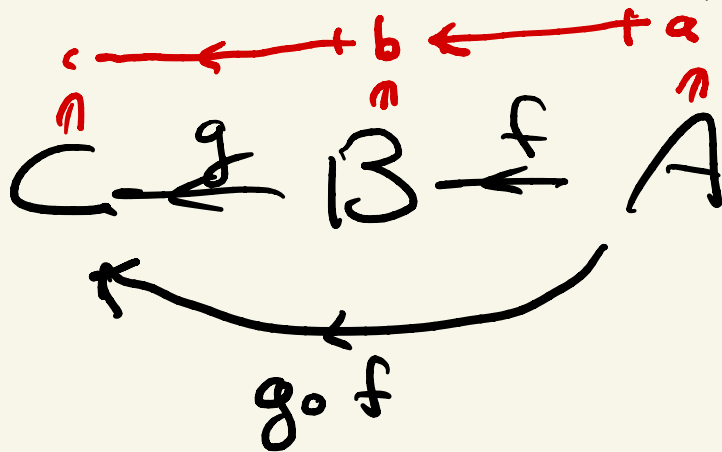


Note: g is one to one if every horizontal line hits the graph at most once.

f is surjective if every horizontal line hits the graph at least once.

Sketch:

U: Need



L: $\forall c \in C$ want $\exists b \in B$
with $g(b) = c$.

I: Know: $g \circ f$ is onto so
 $\exists a \in A$ with $(g \circ f)a = c$


so take $b = f(a)$ in B ✓

Sketch:

L: $P \not\Rightarrow Q$

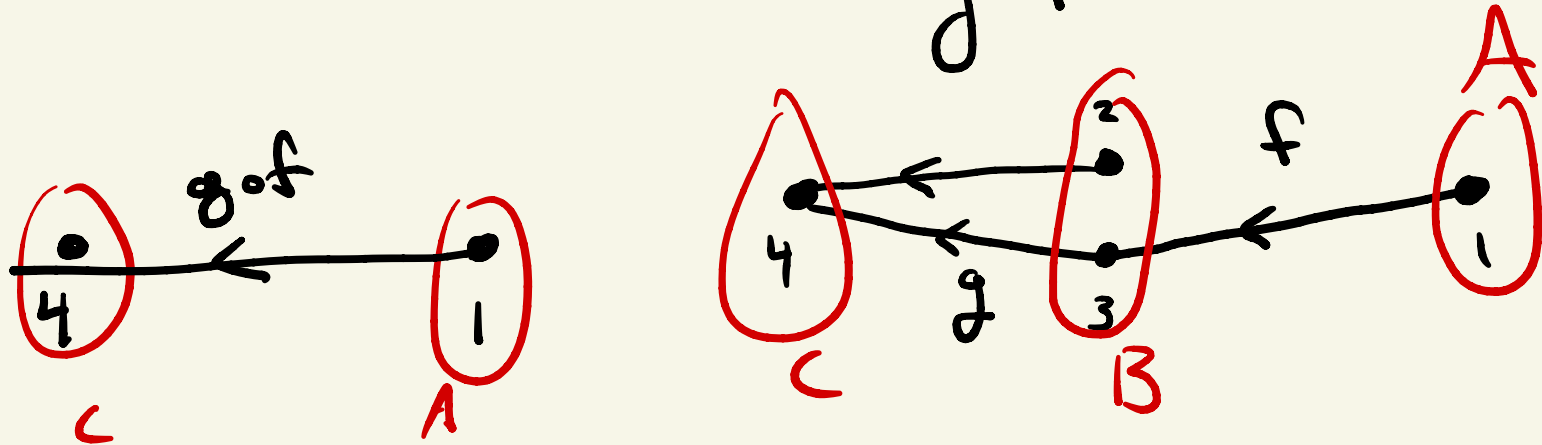
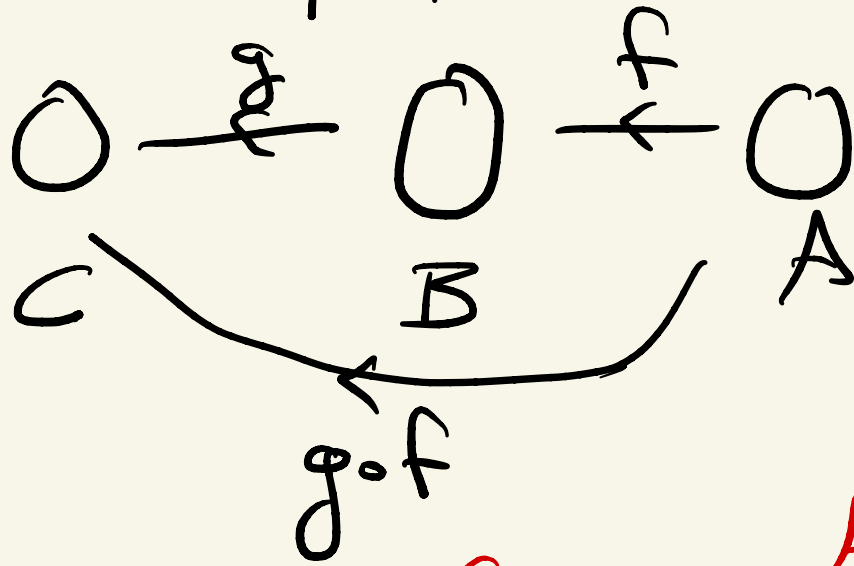
$\rightarrow P$

Q is f is onto

equivalent:

$\exists f, g$ as above with f not onto.

I: Find an example:



If $A \subseteq B$

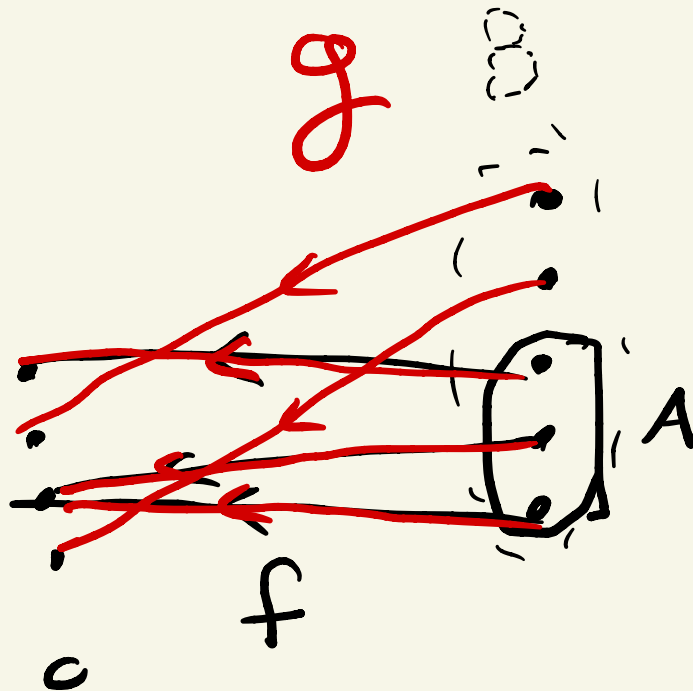
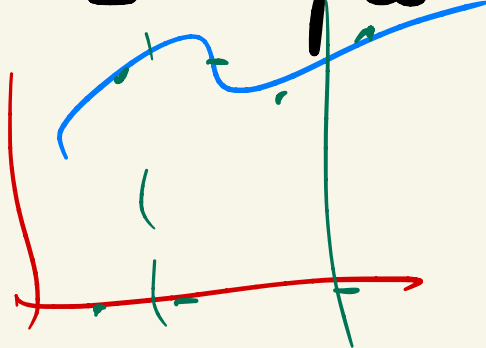
restriction! If $f: B \rightarrow C$
then the restriction of f to A

is $f|_A: A \rightarrow C$

with $f|_A(a) = f(a)$

extension: f is an extension of $f|_A$.

Example:



f is a fn
an ext. from A to C
 g to B of f

Cor 4.4.3:

If a relation f from A to B
is a bijective function
then the relation f^{-1} from B to A
is also a bijective function.

Proof: Assume f is a relation
from A to B which is a
bijective function.

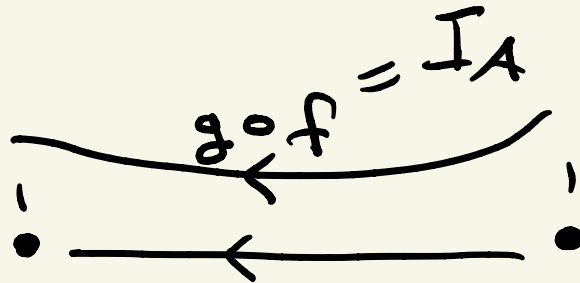
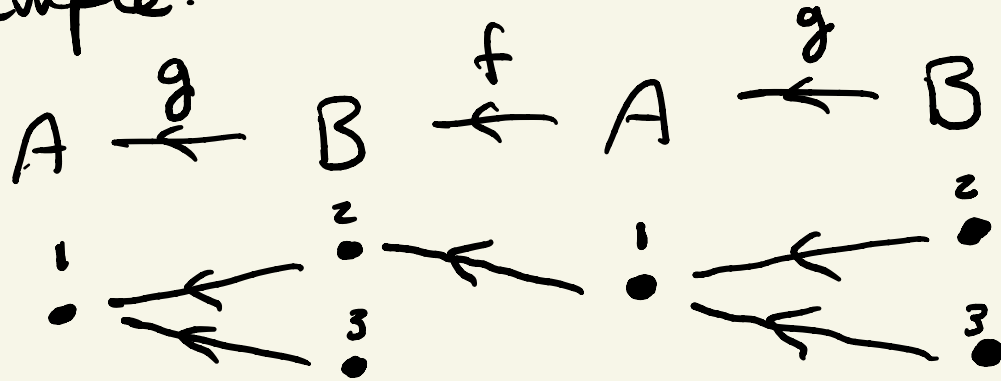
Since f is bijective if b is in B

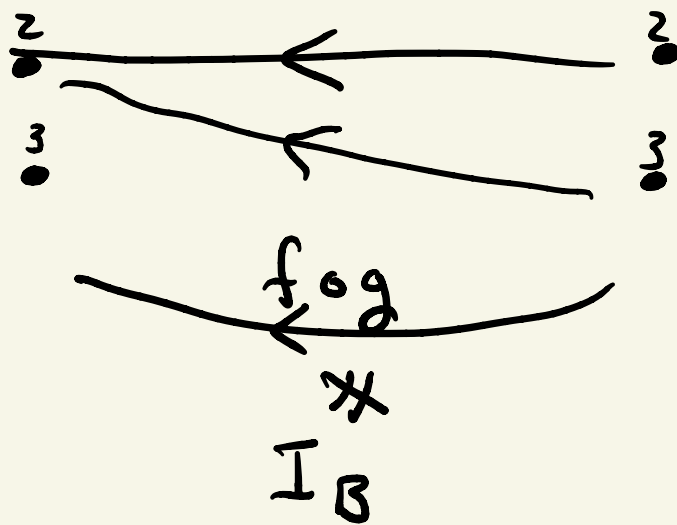
there is a unique a in A with
 (a, b) in f or equivalently (b, a) in f^{-1}
so f^{-1} is a function.

Since f is a function if a is in A
there is a unique b in B with
 (a, b) in f or equivalently (b, a) in f^{-1}
so f^{-1} is a bijection. qed.

Sketch for Claim:

Existential so look for an example.





Sketch for Thm 4.4.4.a:

L : iff so need $\Rightarrow$ $\Leftarrow$] two things,

$\Rightarrow$ Assume $f = g^{-1}$ and

show 4 things:

$$g \circ f \subseteq I_A$$

$$g \circ f \supseteq I_A$$

$$f \circ g \subseteq I_B$$

$$f \circ g \supseteq I_B$$

$\Leftarrow$ Assume $g \circ f = I_A$ and $f \circ g = I_B$
and show 2 things:

$$f \subseteq g^{-1}$$

$$f \supseteq g^{-1}$$

I: $\Rightarrow$ If $(a, \tilde{a}) \in g \circ f$
then $\exists b \in B$ with
 $(a, b) \in f$ and $(b, \tilde{a}) \in g$.

hence ~~$(\tilde{a}, b) \in f = g^{-1}$~~

~~and since f is a fn. $a = \tilde{a}$.~~

hence $(b, a) \in g = f^{-1}$

and since g is a fn.
have $a = \tilde{a} = g(b)$,

therefore $g \circ f \subseteq IA$

Claim: $f: \mathbb{Z}_5 \rightarrow \mathbb{Z}_5$

① with $f(\bar{x}) = (\bar{x})^2$ is not a bijection

Claim: $f: \mathbb{Z}_5 \rightarrow \mathbb{Z}_5$

② with $f(\bar{x}) = (\bar{x})^3$ is a bijection.

Claim: $f: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$

③ with $f(a, b) = 2^{a-1}(2b-1)$
is a bijection.

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Proof of Claim ③:

f is onto since if n is a natural number
then $n = 2^2 \beta$ with β odd, $\alpha \geq 0$ and $\beta \geq 1$.

Take $a = \alpha + 1$ and $b = \frac{\beta + 1}{2}$. Hence $f((a, b)) =$

$$2^a (2^{\frac{\beta + 1}{2}} - 1) = n$$

f is 1-1 since if $2^a (2b - 1) = 2^{\tilde{a}} (2\tilde{b} - 1)$

and if (for contradiction) $a > \tilde{a}$

then $2^{a - \tilde{a}} (2b - 1) = 2\tilde{b} - 1$ is both even

and odd (a contradiction) so $a \leq \tilde{a}$.

similarly $\tilde{a} \leq a$ so $a = \tilde{a}$.

Hence $2^{a-1}(2b-1) = 2^{a-1}(2\tilde{b}-1)$

so $2b-1 = 2\tilde{b}-1$

so $b = \tilde{b}$. qed.

Example: Some bijections:

① $\alpha: (0,1) \rightarrow (0,3)$

with $\alpha(x) = 3x$ is a bijection.

Check: since if $g(y) = \frac{1}{3}y$

$$\text{check: } (a \circ g)(y) = a(g(y)) = 3 \cdot \frac{1}{3} y = y.$$

$$\text{so } a \circ g = I_{(0,3)}$$

$$\text{and } (g \circ a)(x) = g(a(x)) = \frac{1}{3} \cdot 3 \cdot x = x$$

$$\text{so } g \circ a = I_{(0,1)}$$

Find a bijection

$$f: (0, 1] \rightarrow (0, 1)$$

$$\text{and find } f\left(\frac{1}{2}\right) \text{ and } f\left(\frac{1}{\pi}\right)$$

and $f(1)$.

Using: $b: (0,1) \rightarrow (1, \infty)$ is a bijection.
 $b(x) = \frac{1}{x}$

and $c: (0,1] \rightarrow [1, \infty)$ is a biject.
 $c(x) = \frac{1}{x}$

and $d: \mathbb{W} = \{0,1,2, \dots\} \rightarrow \mathbb{N}$ is a biject
 $d(n) = n+1$

and

$$e: [1, \infty) \rightarrow (1, \infty)$$

$$e(x) = \begin{cases} x+1 & \text{if } x \in \mathbb{N} \\ x & \text{otherwise} \end{cases}$$

is a bij.

Some counting: If $\bar{A} = 3$ find:

(a) $\overline{A \times A}$

(b) $\overline{P(A)}$

(c) $\overline{P(A \times A)}$

(d) $\{f \mid f \text{ is a relation from } A \text{ to } A\}$

(e) $\{f \mid f \text{ is an equivalence reln on } A\}$

(f) $\{f \mid f \text{ is a function from } A \text{ to } \mathbb{N}_2\}$
 $\parallel$
 $\{1, 2\}$

$$f \subseteq A \times A$$

equiv to partitions

- ① $\{f \mid f \text{ is a function from } A \text{ to } A\}$
 ② $\{f \mid f \text{ is a bijection from } A \text{ to } A\}$

Ans: ① $\overline{A \times A} = \overline{A} \cdot \overline{A} = 3 \cdot 3 = 9$

② $\overline{P_A} = 2^{\overline{A}} = 2^3 = 8$

③ $\overline{\overline{P(A \times A)}} = 2^9 = 512$

④ the same as ③ since
 a subset of $A \times A$ is the
 same as a reln from A to A .
 so 512 possible relations.

(e) the same as partitions of $\{1, 2, 3\}$
 $= \{ \{ \{1, 2, 3\} \}, \{ \{1, 2\}, \{3\} \}, \{ \{1, 3\}, \{2\} \}, \{ \{2, 3\}, \{1\} \}, \{ \{1\}, \{2\}, \{3\} \} \}$

$$= 5$$

(f) $\overline{N_2} \overline{A} = 2^3 = 8$

(g) $\overline{A} \overline{A} = 3^3 = 27$

(h) $6 = \overline{A}!$

- 3 choices
- 2 choices
- 1 choice

Thm: 5.1.1

$\approx$ is an equivalence relation
on the collection of sets.

Proof: Need to check:

→ symmetry.

- reflexivity.

- transitivity.

Check: symmetry: Need: If $A \approx B$ then $B \approx A$.

That is if $\exists f: A \rightarrow B$ a bij. function
then $\exists g: B \rightarrow A$ a bij. function

Proof: Assume $f: A \rightarrow B$ is a bij.
 function. Hence $f^{-1}: B \rightarrow A$ is also a
 bij. function. *qed.*

Overview of small cardinalities:
 no two columns are equivalent.

$\overline{\overline{\mathbb{P}\mathbb{N}}}$	$\not\approx \mathbb{Q} \approx \mathbb{Q}_{>0} \approx \mathbb{N} \times \mathbb{N}$	$(0,1]$	$\not\approx \overline{\overline{\mathbb{P}(\mathbb{N})}}$
	$\mathbb{Z}$	$\mathbb{S}$	
	$\mathbb{Z}$	$\mathbb{R}$	
	$\mathbb{S}$	$\mathbb{S}$	
	$\mathbb{N}$	(a,b)	
	$\mathbb{S}$	$\mathbb{S}$	
$\emptyset$	$\mathbb{N}_1 \not\approx \mathbb{N}_2 \dots \not\approx \mathbb{N}_n \not\approx$	$\not\approx \mathbb{N} \not\approx$	$(0,1)$

A

0 1 2 ... n ... $\aleph_0$ $\subset$

name of
A

Thm: If $r < n$ then $N_r \neq N_n$.

Proof strategy:

Show there is no $N_n \xrightarrow{f} N_r$
which is injective (so no bijection)

Use induction. with $P(n)$ the proposition
that $\forall r < n \quad \nexists f: N_n \rightarrow N_r$
which is injective.

Possible base steps for induction:

$P(1)$ means: there is no injective map
from $\{1\}$ to $\emptyset$

which is true since there is
no function from $\{1\}$ to $\emptyset$.

$P(2)$ means: there is no injective
map from $\{1, 2\}$ to $\{1\}$ or $\emptyset$
 $r=1$ $r=0$

but there is only one fn $f: \{1, 2\} \rightarrow \{1\}$
and $f(1) = f(2) = 1$ so not inj.

and there are no ~~any~~ fns from $\{1, 2\}$ to $\emptyset$.

210310

Thm 5.2.1: If $n \in \mathbb{N}$ then
 $\mathbb{N}_n \not\approx \mathbb{N}$. ($\mathbb{N}$ is infinite),

Proof Plan: Show $\nexists f: \mathbb{N} \rightarrow \mathbb{N}_n$

injective.

Use 5.1.9: $\nexists g: \mathbb{N}_{n+1} \rightarrow \mathbb{N}_n$ injective.

Use contradiction.

Assume $f: \mathbb{N} \rightarrow \mathbb{N}_n$ is injective.

Hence $f|_{\mathbb{N}_{n+1}}: \mathbb{N}_{n+1} \rightarrow \mathbb{N}_n$ is injective.

a contradiction.

q.e.d.

Proof of 5.2.4:

Assume $f: \mathbb{N} \rightarrow (0,1)$.

Write a_n for the n^{th} decimal digit of $f(n)$

and choose $b_n \neq a_n$ and $B \in (0,1)$

with digits b_n .

Claim: $B \notin \text{Im}(f)$ so f is not onto.

Pf of claim [For contradiction assume $f(m) = B$
so $a_m = b_m$ a contradiction.

q.e.d.

Proof of 5.4.3:

Assume $f: A \rightarrow \mathcal{P}(A)$

Write $B = \{a \in A \mid a \notin f(a)\}$

Claim $B \notin \text{Im}(f)$ so f is not onto.

Pf of claim:

For contradiction assume $f(a) = B$

Two cases: ① $a \in B$
or ② $a \notin B$.

In case ① $a \in f(a)$ so $a \notin B = f(a)$
a contradiction.

In case ② $a \notin f(a)$ so $a \in B = f(a)$
another contradiction. q.e.d.

210312) (Reread §1.7)

Proofs: Methods:

direct.

cases

contrapositive (not the converse).

contradiction

induction

(PMI, Generalized).

Possible hints from logic:

$\hookrightarrow$

start with suppose

maybe direct

$\exists$

.. ..

choose

maybe construct

iff $\Leftrightarrow$ two things $\Rightarrow$ $\{ \Leftarrow$
= sets " " $\subseteq$ $\{$ $?$

Office hrs at 1 on Monday.

Ans: Outline @:

Need to show: reflexive (r)
symmetric (s)
transitive (t).

Idea:

$$(r): \quad \forall x \in \mathbb{R} \quad (x, x) \in S$$

since $x - x = 0 \in \mathbb{Q}$ ✓

$$(s): \quad \forall x, y \in \mathbb{R} \quad (x, y) \in S \Rightarrow (y, x) \in S$$

Assume $x - y \in \mathbb{Q}$

hence $-(x - y) = y - x \in \mathbb{Q}$ ✓

$$(t) \quad \forall x, y, z \in \mathbb{R} \quad ((x, y) \in S \wedge (y, z) \in S) \Rightarrow (x, z) \in S.$$

Assume $x - y \in \mathbb{Q}$ and $y - z \in \mathbb{Q}$

hence $x - y + y - z = x - z \in \mathbb{Q}$ ✓

© Show $\mathbb{R}$ is denumerable:

Proof idea: Need $\overline{\mathbb{N}} \approx \mathbb{N}$

or equiv: $\exists \pi \xrightarrow{f} \mathbb{N}$ a bijection.

Note also know: (denumerable means $\overline{A} = \aleph_0$),
 $\{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$

$\mathbb{Z}, \mathbb{N}, \mathbb{N} \times \mathbb{N}, \mathbb{Q}$

so it suffices to find a bijectn.
from $\overline{\pi}$ to any of

Consider $\overline{\pi} \xrightarrow{f} \mathbb{Q}$

with $f(x) = x - \pi$

(Proved last time:
 $\exists \mathbb{N} \xrightarrow{f} \mathbb{R}$
with f onto
so $\overline{\mathbb{N}} \neq \overline{\mathbb{R}}$)

since $x \in \overline{\pi}$ know $(\pi, x) \in S$

so $\pi - x \in \mathbb{Q}$.

so $x - \pi \in \mathbb{Q}$.

$$\text{eg } f\left(\underbrace{\pi + \frac{2}{3}}_{\in \overline{\pi}}\right) = \frac{2}{3}$$

$\pi + \frac{2}{3} - \pi$

Finally: show f is a bijection:

How to do this: $\left[\text{Find an inverse to } f. \right.$

or [show onto
and show 1-1

✓.

Consider $g(y) = y + \pi$

check: $(f \circ g)(y) = f(g(y)) = f(y + \pi) = y$

and $(g \circ f)(x) = g(f(x)) = g(x - \pi) = x$
q.e.d.

Direct Proof of onto

Need: $\forall y \in \mathbb{Q} \quad \exists x \in \pi$ with $f(x) = y$.

Pf: Assume $y \in \mathbb{Q}$ take $x = y + \pi$

so $f(x) = f(y + \pi) = y + \pi - \pi = y$

$$f(x) = x - \pi$$

and $x \in \overline{\pi}$ since

$$\pi - x = \pi - (y + \pi) = -y \in \mathbb{Q}.$$

$$\overline{\pi} = \{x \mid \pi - x \in \mathbb{Q}\}$$

