## Notes on Rasmussen's Proof of Milnor's Conjecture

From last time we have
Definition. A cobordism between links $L_{0}, L_{1}$ is a smooth surface $W \subseteq S^{3} \times[0,1]$ with $\left.W\right|_{S^{3} \times\{0\}}=$ $L_{0}$ and $\left.W\right|_{S^{3} \times\{1\}}=L_{1}$ and boundary $\partial W=L_{0} \cup L_{1}$.

We also recall the following:

- We may decompose a cobordism $W$ into a series of elementary cobordisms corresponding to Reidemeister moves and Morse moves.
- A cobordism $W$ induces a map $\Phi_{W}: H_{\text {Lee }}\left(L_{0}\right) \rightarrow H_{\text {Lee }}\left(L_{1}\right)$
- Two cobordisms are equivalent if we can relate them by a series of "movie moves". Theorem from last week: $\Phi_{W}$ is well-defined.

We will show that for knots $\Phi_{W}$ is an isomorphism, but first we introduce a related notion:
Definition. Two links $L_{0}, L_{1}$ are concordant if there exists an embedding $f: L_{0} \times[0,1] \rightarrow S^{3} \times[0,1]$ such that $f\left(L_{0} \times\{i\}\right)=L_{i} \times\{i\}$ for $i \in[0,1]$. A knot $K$ is called slice if $K$ is concordant to the unknot.

More intuitively: Imprecisely: equivalent to saying that there exists a cobordism of genus 0 between $L_{0}$ and $L_{1}$ or $L_{0}$ and $L_{1}$ are connected by a cylinder in $S^{3} \times[0,1]$ with boundary components $L_{i}$. With this formulation, a knot $K$ is slice if it bounds a disk in $B^{4}$ (see example).
Fact: Concordance gives an equivalence relation on knots by setting $K_{1} \sim K_{2}$ if $K_{1} \# \overline{K_{2}}$ is slice. The set of knots modulo this relation forms an abelian group with the connect sum operation. This group is called the concordance group of knots in $S^{3}$ and is denoted $\operatorname{Conc}\left(S^{3}\right)$.

Definition. The slice genus of a knot, $K$ is the minimum genus $g$ of a connected, orientable surface $S$ (smoothly) embedded in $B^{4}$ that has $K$ as a boundary.

In general, slice genus is difficult to compute. Lisa Piccirillo's proof that the Conway knot is not slice completed the classification of slice knots under 13 crossings only 2 years ago.

The following theorem provides a lower bound on slice genus
Theorem 1.

$$
|s(K)| \leq 2 g_{*}(K) \leq 2 g(K)
$$

where $g_{*}(K)$ denotes the slice genus.
To prove this theorem, we need to examine how the canonical generators of $H_{\text {Lee }}$ behave under cobordism.

Proposition. Let $W$ be a weakly connected cobordism from $L_{0}$ to $L_{1}$. Then $\Phi_{W}\left(z_{o_{0}}\right)$ is a nonzero multiple of $z_{o_{1}}$ where $z_{o_{i}}$ denotes a canonical generator of $H_{L e e}\left(L_{i}\right)$ labeled by orientation $o$ on $L_{i}$.
Definition. A cobordism $W$ is weakly connected if all components of $W$ have a boundary component in $L_{0}$.

The following corollary follows immediately from this proposition:
Corollary. If $W$ is a connected cobordism between knots, then $\Phi_{W}$ is an isomorphism.

Proof. (of proposition) A cobordism $W$ between two links $L_{1}$ and $L_{2}$ induces a morphism $\Phi_{W}$ on the Lee homology as follows:

For $W=W_{1} \cup \cdots \cup W_{k}$, with corresponding links $L_{i}$, define $\Phi_{W}=\Phi_{W_{1}} \circ \cdots \circ \Phi_{W_{k}}$, where for each $i, \Phi_{W_{i}}: H_{\text {Lee }}\left(L_{i}\right) \rightarrow H_{\text {Lee }}\left(L_{i+1}\right)$. We can then prove the proposition by induction on $i$. If $W_{i}$ is an elementary cobordism corresponding to a Reidemeister move, we have (from our earlier discussion of Lee homology) a filtered map (of $q$-degree zero) $\rho_{*}$ sending canonical generators to canonical generators. Therefore, it suffices to check that the proposition holds for an elementary cobordism corresponding to attaching a handle. For such $W_{i}$ we apply $\iota^{\prime}, m^{\prime} / \Delta^{\prime}$, or $\epsilon^{\prime}$ at each vertex of our cube of resolutions for 0,1 , and 2 -handle moves respectively. Whether we apply $m^{\prime}$ or $\Delta^{\prime}$ depends on whether the 1 -handle merges circles or splits circles at each vertex of the cube of resolutions (see example).

Divide the components of $W_{i}$ into two sets:
Components of the 1st type have a boundary component in $L_{0}$
Components of the 2 nd type have no boundary component in $L_{0}$
We call an orientation $o$ on $W_{i}$ permissible if it agrees with the orientation on $S$ on components of the first type.

We will prove the slightly stronger statement that

$$
\Phi_{W_{i}}\left(z_{o}\right)=\sum_{I} a_{I} z_{o_{I}}
$$

for permissible orientations $I$ of $W_{i+1}$ arguing by induction on $i$.

- 0-handle: The generators of $H_{\text {Lee }}$ of $L_{i}$ are given by $z_{o} \otimes a, z_{o} \otimes b, z_{\bar{\sigma}} \otimes a$ and $z_{\bar{o}} \otimes b$. The $\operatorname{map} \Phi_{W}\left(z_{o}\right)=\iota\left(z_{0}\right)=z_{0} \otimes \frac{1}{2}(a-b)$ is then a nonzero multiple of the two generators with the appropriate orientation.
- 1-handle: If $O_{I}$ agrees with $o_{i}$ then the two strands involved in the 1-handle move have opposite orientations, hence are both labeled $a$ or both labeled $b$. We have

$$
\begin{array}{rr}
m(a \otimes a)=2 a & \Delta(a)=a \otimes a \\
m(b \otimes b)=-2 b & \Delta(b)=b \otimes b
\end{array}
$$

So in this case, $\Phi_{W_{i}}\left(z_{o_{i}}\right)$ is a nonzero multiple of $z_{o_{I}}$.
If $o_{I}$ is not compatible with $o_{i}$ then we have two strands pointing in the same direction with different labels, so $\Phi_{W_{i}}\left(z_{o_{i}}\right)=0$ because $m(a \otimes b)=m(b \otimes a)=0$. From our discussion of Lee homology before, we know that we cannot split a circle into two circles with opposite labels, so we do not need to consider $\Delta^{\prime}$ for this subcase.

To complete the argument, we need to examine how the orientations behave under $\Phi_{W_{i}}$. If attaching the 1 -handle adds a component to $L_{i}$, then the number and type of components is the same, so the set of permissible orientations is preserved (naturally identified) If we merge two components, then we have three cases:

- The merge only involves one component of $W_{i}$
- The merge involves two components of $W_{i}$ of the first type
- The merge involves two components, one of which is of the second type

In the first two cases, the set of permissible orientations is preserved, as above. In the third case, the set of permissible orientations is halved, but we still preserve canonical generators.

- 2-handle: an orientation on $W_{i}$ uniquely extends to an orientation on $W_{i+1}$. Then $\epsilon(a)=$ $\epsilon(b)=1$ implies that $\Phi_{W}\left(z_{o_{i}}\right)=z_{o_{i+1}}$
Note that here, we might have two different orientations on $W_{i+1}$ inducing two different orientations on $L_{i+1}$. But this can only happen when we have our two-handle capping a closed component and this violates the weakly connected hypothesis.
(See example)
Important note: We can compute the filtered $q$-degree of elementary cobordisms corresponding to handle attaching maps using our formula

$$
q \text { grading }=\operatorname{degree}_{A}-|v|-\# \text { circles }-n_{+}-2 n_{-} .
$$

In our example, everything is preserved, except $n_{+}$, which decreases by 1 . This holds in general and the degree of a 1 -handle attaching is 1 .

For 0-handle maps, we add a circle and everything else is preserved, so the $q$ grading is -1 . For 2 -handle maps, we add a delete a circle, but degree ${ }_{A}$ decreases by 2 because we remove an $x$ from the tensor product, so the $q$ grading is -1 again. The filtered degree of an elementary cobordism corresponding to a handle attaching move is then the (negative) Euler characteristic, so the degree of the map $\Phi_{W}$ with our sign convention is

$$
-\chi(W)=-(2-2 g-2)=2 g=-\# 0 \text {-handles }+\# 1 \text {-handles }-\# 2 \text {-handles. }
$$

Now we wish to use $\Phi_{W}$ to get a relationship between $s(K)$ and slice genus of $K$ in order to prove theorem 1.

Proof. (of Theorem 1) Let $K$ be a knot with slice genus $g_{*}(K)=g$. Consider $x \neq 0$ with minimal grading in $H_{\text {Lee }}(K)$. We then have a cobordism $W$ of genus $g$ from $K$ to the unknot $U$. By our corollary above, the induced map $\Phi_{W}$ is an isomorphism with $s_{\min }\left(\Phi_{W}(x)\right) \geq-1=s_{\text {min }}(U)$ (recall that for the unknot, $s(U)=0$ and $s_{\text {min }}(U)=-1$. Since the $q$-degree of $\Phi_{W}$ is also $2 g$, we have $s(x) \geq-2 g-1$ since $\Phi_{W}$ could have added at most $2 g$ to the $q$-degree of $x$. We chose $x$ to have minimal grading, so $s_{\min }(K) \geq-2 g-1$. Hence, $s(K) \geq-2 g$. The bound on $s(K)$ follows from the fact that $s(\bar{K})=-s(K)$

Corollary. For a positive knot $K, s(K)=g_{*}(K)=g(K)$.
Proof. Last time we showed $g(k) \leq s(K)$ since a surface given by Seifert's algorithm has genus $s(K)$. With the inequality from the theorem above, we get the desired equality

Corollary. (Milnor conjecture) The slice genus of a $(p, q)$-torus knot is $\frac{(p-1)(q-1)}{2}$.
Proof. The genus of the $(p, q)$-torus knot is known to be $\frac{(p-1)(q-1)}{2}$. Since torus knots are positive, the above corollary implies that $g_{*}(K)=g(K)$.

We also have a nice algebraic property of the $s$-invariant. It behaves nicely with respect to connect sum and we have the following theorem:

Theorem 2. The map $s$ induces a homomorphism from $\operatorname{Conc}\left(S^{3}\right)$ to $\mathbb{Z}$, where $\operatorname{conc}\left(S^{3}\right)$ denotes the concordance group of knots in $S^{3}$

Proof. See properties of $s(K)$ under connect sum for additivity (basic idea is that canonical generator of connect sum is tensor product of canonical generators of knots). To show that $s(K)$ is a well defined map, consider concordant knots $K_{1}$ and $K_{2}$. Then the connect sum is slice, so $0=s\left(K_{1} \# \overline{K_{2}}\right)=s\left(K_{1}\right)-s\left(K_{2}\right)$ and thus, $s\left(K_{1}\right)=s\left(K_{2}\right)$.
( Cor: Knots differing by a single crossing satisfy $s\left(K_{-}\right) \leq s\left(K_{+}\right) \leq s\left(K_{-}\right)+1$ )

